

Intro to Crypto

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Agenda – Intro To Crypto

- Some background
- Crypto Functions
 - Symmetric algorithms
 - Asymmetric algorithms
 - Hashes
 - PIN Support



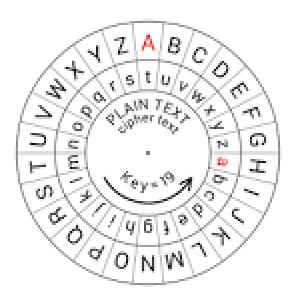


Historical Ciphers



Scytale





Caesar Cipher, Key = 3

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PDLQIUDPH

Vigenere Square, Key = BOYD

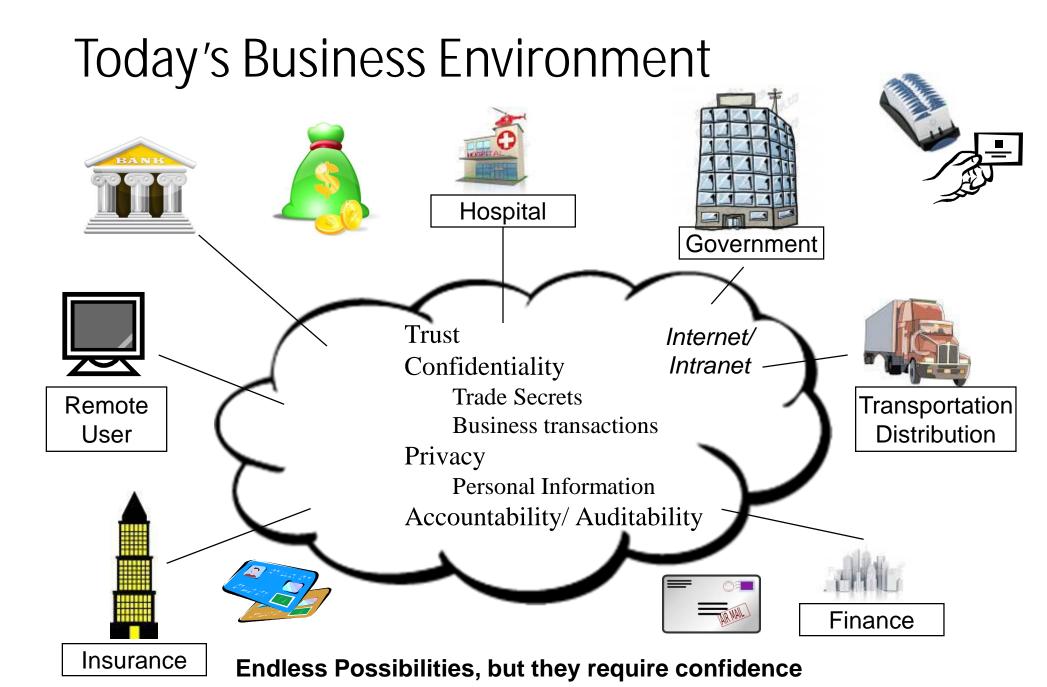
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BOYDBOYDB

NOGQGFYPF











What is Cryptography?

Cryptography (or cryptology; from <u>Greek</u> κρυπτός, *kryptos*, "hidden, secret"; and γράφω, *gráphō*, "I write", or - λογία, <u>-logia</u>, respectively)[1] is the practice and study of hiding <u>information</u>. In modern times cryptography is considered a branch of both <u>mathematics</u> and <u>computer science</u> and is affiliated closely with <u>information theory</u>, <u>computer security</u> and <u>engineering</u>.

From Wikipedia





Cryptographic Functions

Data Confidentiality

Symmetric – DES/TDES, AES

Asymmetric – RSA, Diffie-Hellman, ECC

Data Integrity

Modification Detection Message Authentication Non-repudiation

Financial Functions Key Security & Integrity

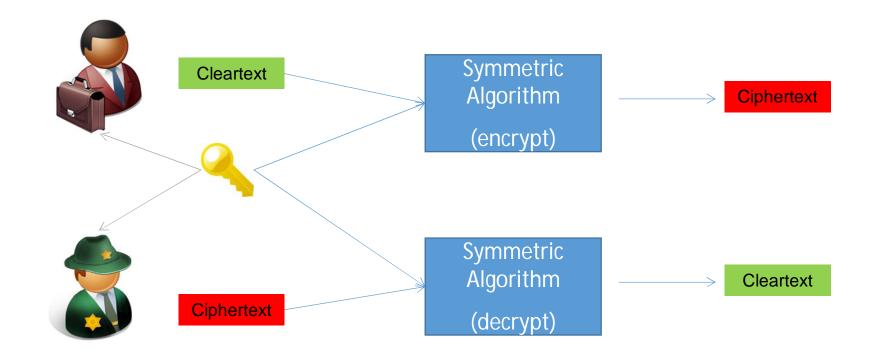






Confidentiality – Symmetric Algorithms

Symmetric - One key shared by both parties







Symmetric Algorithms

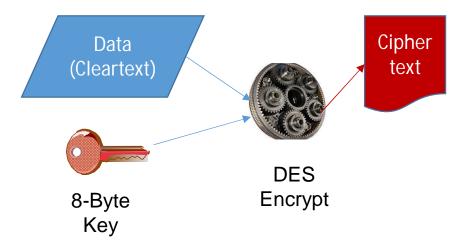
- Symmetric
 - DES/TDES*
 - AES*
 - Blowfish / Twofish
 - Serpent
 - IDEA
 - RC2 / RC4
 - Skipjack

*Supported on IBM Hardware





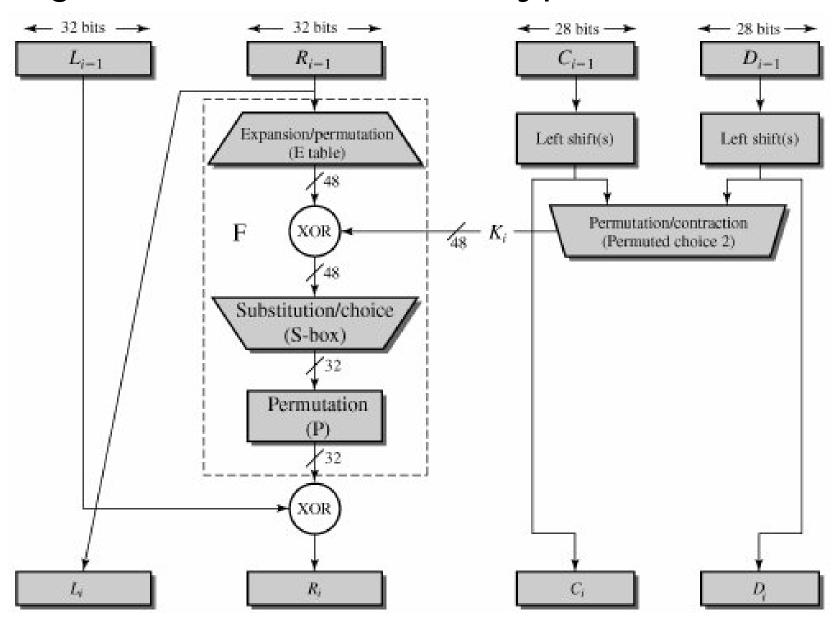
DES Algorithm - Encrypt







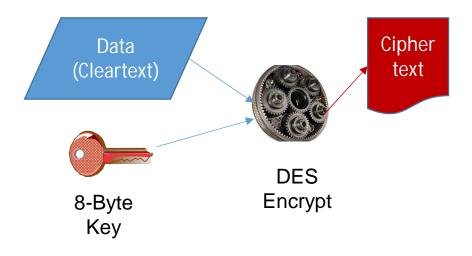
Single Round of DES Encrypt

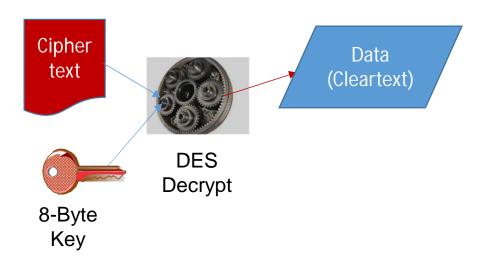






DES Algorithm - Decrypt

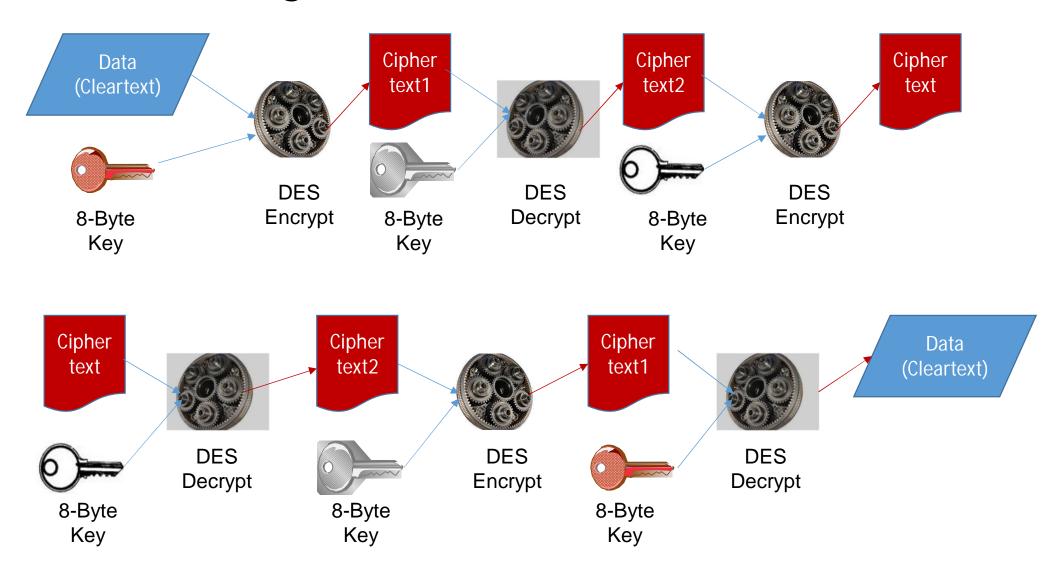








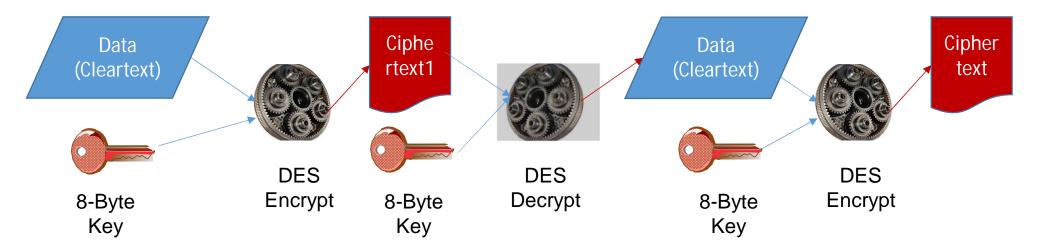
TDES Algorithm







TDES Algorithm







Data Confidentiality - AES

- Rijndael Algorithm
 - Block Cipher (16-byte blocks)
 - 128-, 192-, 256-bit key length
 - 128 bit key=> 3.4x10**38 (340 Undecillion)
 - 192 bit key=> 6.2x10**57 (6.2 Octodecillion)
 - 256 bit key=> 1.1x10**77 (almost a Googol)
 - Multiple round
 - Four steps per round (Byte substitution, shift row, mix column, add round key)

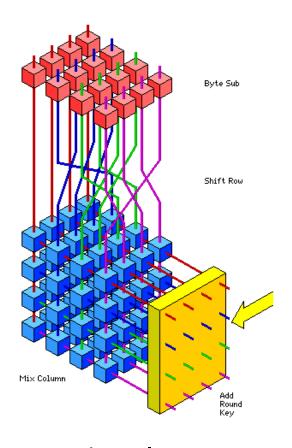


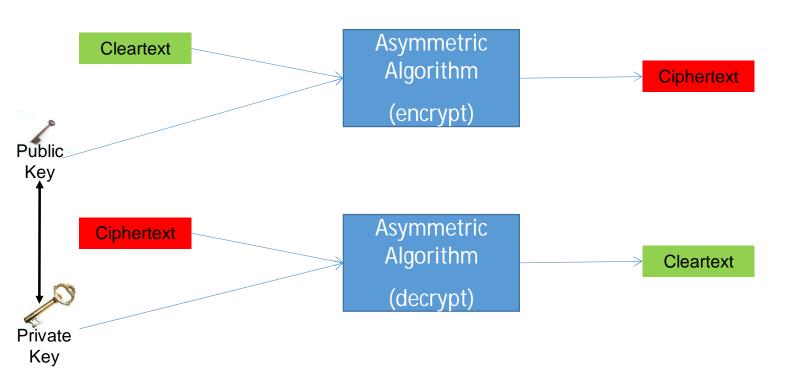
Image from http://www.esat.kuleuven.ac.be/~rijmen/rijndael





Secrecy Algorithms - Asymmetric

 Asymmetric – two different, but mathematically related keys (public and private)

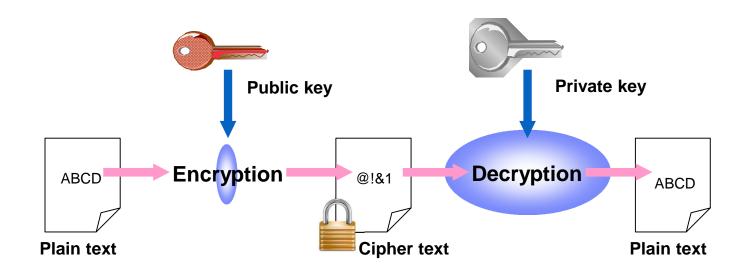






Asymmetric Algorithms

- Public Key Architecture PKA
 - RSA factorization
 - Diffie-Hellman logarithmic
 - Elliptic Curve point multiplication







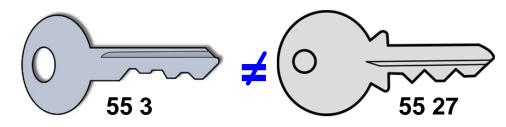
Generating RSA Keys

- RSA Keys consists of two parts, a modulus (N) and an exponent (E for the public key; D for the private key)
 - Public Key => N E
 - Private Key => N D
- The modulus is calculated by multiplying two prime numbers (P & Q) together
 - P = 5 Q = 11 (prime numbers and should be very large)
 - $N = P \times Q => 5 \times 11 = 55$
- Next, select an odd number, E, that will be the exponent for the public key
 - Good values include 3 or 65537 (64K+1) or 5, 17 or 257 with HCR77C0

• Finally, calculate the exponent for the private key, D, where

$$1 = (D * E) MOD ((P-1)(Q-1)) => 1 = (D * 3) MOD ((5-1)(11-1))$$

• In our example, solve for 1 = (D * 3) MOD 40 => D = 27!







Encipher the Message 'MFC'

Public Key (N E) => 55 3

Private Key (N D) => 5527

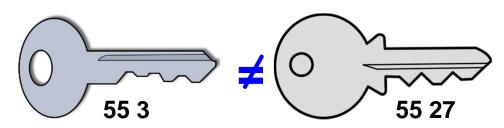
Convert characters to numeric (a=1, b=2, c=3, etc.)

$$'M' = 13; 'F' = 6; 'C' = 3;$$

ciphertext = (cleartext**E) Mod N

- For 'M' (13**3) MOD 55 => 2197 MOD 55 = 52
- For 'F' (6**3) MOD 55 => 216 MOD 55 = 51
- For 'C' (3**3) MOD 55 => 9 MOD 55 = 27

Ciphertext is 52 51 27





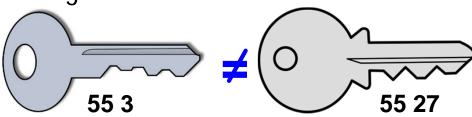


Decipher the message 52 51 27

Public Key (N E) \Rightarrow 55 3

Private Key (N D) => 55 27 Cleartext = (ciphertext**D) MOD N

- For 52 52**27 MOD 55 = 13 (52**27 = 2.1482769967144679013436706816572e+46)
- For 51 51**27 MOD 55 = 6 (51**27 = 1.2717295264013893903823981998699e+46)
- For 27 27**27 mod 55 = 3 (27**27 = 4.4342648824303776994824963061915e+38)
- My decrypted message is 13 6 3 => "M" "F" "C"







ECC Algorithm

Effective Key Size (bits)				
Symmetric	RSA	ECC	J-R	P (-2.35, -1.86) Q (-0.1, 0.836) -R (3.89, 5.62)
80	1024	163		R (3.89, -5.62)
112	2048	224	$\begin{cases} & 1 \\ & & 1 \\ & & 1 \end{cases} \rightarrow x$	P+Q=R=(3.89, -5.62).
128	3072	256	\	
192	7680	384	\;	
256	15360	512	R	
From NIST SP 800-57 Part 1 (Table 2) at www.nist.gov				

Image from crypto.stackexchange.com





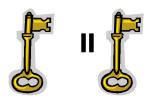
Why Asymmetric and Symmetric Keys?

- Asymmetric
 - plus its strength, can be used to establish a secret between two parties
 - minus expensive in terms of performance



- plus less resource intensive
- minus requires key to be shared securely

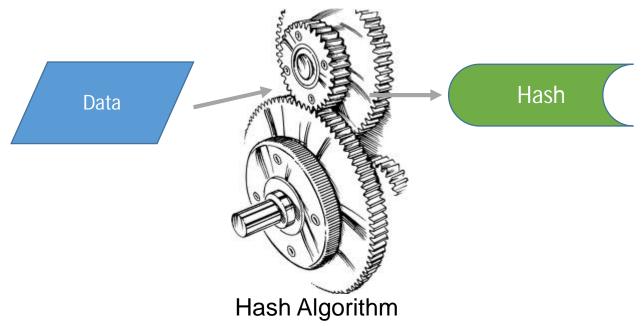








Hashing



- Characteristics of a good hash algorithm
 - One-way can't recover the data from the hash
 - Hard to find collisions
 - The result does not reveal information about the input



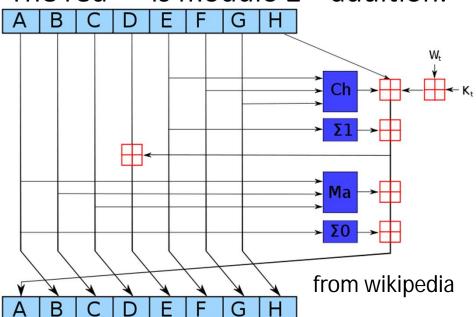


Hashing

 One iteration in a SHA-2 family compression function. The blue components perform the following operations:

$$\begin{array}{l} \operatorname{Ch}(E,F,G) = (E \wedge F) \oplus (\neg E \wedge G) \\ \operatorname{Ma}(A,B,C) = (A \wedge B) \oplus (A \wedge C) \oplus (B \wedge C) \\ \Sigma_0(A) = (A \ggg 2) \oplus (A \ggg 13) \oplus (A \ggg 22) \\ \Sigma_1(E) = (E \ggg 6) \oplus (E \ggg 11) \oplus (E \ggg 25) \end{array}$$

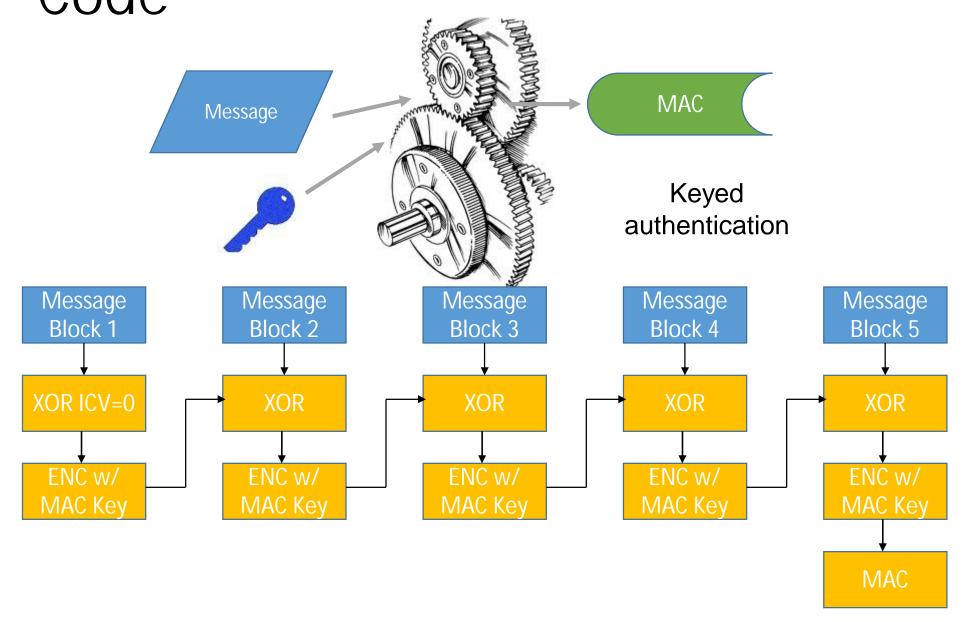
• The bitwise rotation uses different constants for SHA-512. The given numbers are for SHA-256. The red is modulo 2³² addition.







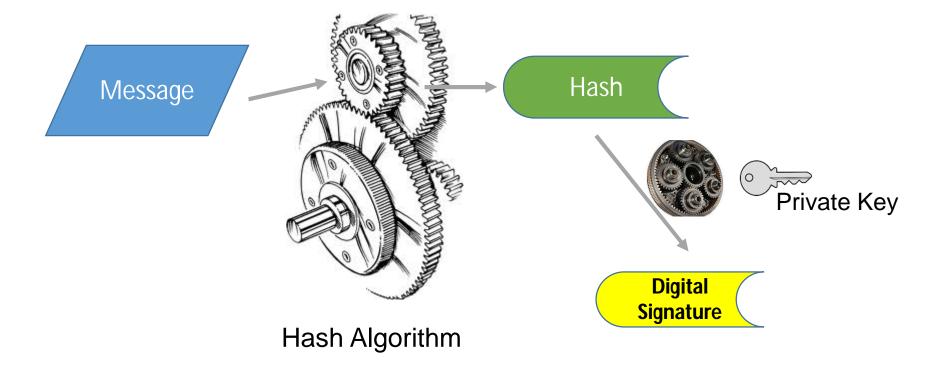
Hashing – Message Authentication Code







Digital Signatures



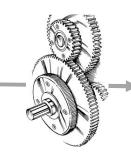






Certificate Request

Version /
Subject Serial Issuer Subject
Name Info Dates Number Algorithms Name Info Public Key



Hash



Certificate
Authority
Private
Key

Digital Signature

Certificate

Version /
Subject Serial Issuer Subject
Name Info Dates Number Algorithms Name Info Public Key





Certificate
Authority
Public Key

Hash





Financial Authentication - PINs

Routing Number:

12345678

Account Number: 9876543210

Sequence Number:

PIN Block:

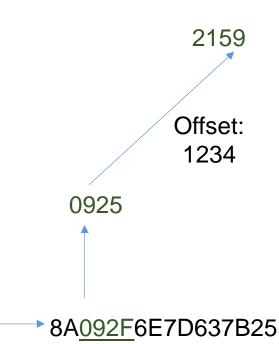
4567898765432101



Pin Block Formats

ECI-2, ECI-3, ISO-0, ISO-1, ISO-2, ISO-3,

VISA-2, VISA-3, VISA-4, 3621, 3624, 4704-FPP



Decimalization Table					
0 -> 0	1 -> 1	2 -> 2	3 -> 3		
4 -> 4	5 -> 5	6 -> 6	7 -> 7		
8 -> 8	9 -> 9	A -> 0	B -> 1		
C -> 2	D -> 3	E -> 4	F -> 5		





References

Cryptography Books

- Bruce Schneier, "Applied Cryptography, Second Edition: Protocols, Algorithms, and Source Code in 'C'", Addison Weley Longman, Inc. 1997
- Simon Singh, "The Code Book", Anchor Books, 1999
- Niels Ferguson, Bruce Schneier, "Practical Cryptography", Wiley Publishing, Inc. 2003

• Free Stuff

 <u>www.schneier.com</u> – Bruce Schneier website, with monthly newsletter Cryptogram







Standards Doc

- RSA
 - PKCS #1 RSA Cryptography Specifications Version 2.2 (https://tools.ietf.org/html/rfc8017)
- ECC
 - https://en.wikipedia.org/wiki/Elliptic-curve_cryptography
 - Also see 'Recommendation for Pair-Wise Key-Establishment Schemes Using Discrete Logarithm Cryptography https://csrc.nist.gov/publications/detail/sp/800-56a/rev-3/final
- AES
 - FIPS 197 Announcing the AES (https://doi.org/10.6028/NIST.FIPS.197)
- DES
 - FIPS 46-3 Data Encryption Standard Withdrawn (http://csrc.nist.gov/publications/fips/fips46-3/fips46-3.pdf)
- TDES
 - SP 800-67 Recommendation for the Triple Data Encryption Algorithm (TDEA) Block Cipher (https://csrc.nist.gov/publications/detail/sp/800-67/rev-2/final)





Questions ...

