Certificates and SSL/TLS: A Look Inside

2021 Edition

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Abstract

Starting with z/OS V2R3, IBM is no longer shipping "standard" certificate-authority (CA) certificates with RACF, putting more responsibility on you to understand and manage certificates on your own. You may have attended other sessions that have shown how to install a certificate in RACF, ACF2 or TSS. This session will give you an understanding of how the certificate process actually works under the covers. It is equally relevant to RACF, ACF2 and TSS systems.

The session will start with a quick review of the underlying technologies and their limitations in the absence of certificates: secret key, public key, and digital signatures and hashes; and go on to cover in detail the protocol flows with server certificates, intermediate certificates, CA certificates and revocation lists. Finally the session will introduce you briefly (with resources for further learning) to advanced topics such as Alternative Names, client certificates, elliptic curve, Diffie-Hellman, code signing, and more.

It is "pure" certificate and SSL/TLS technology. If you are looking to understand what is inside a certificate you are in the right place.

About the Presenter

Charles has been writing mainframe software for longer than he cares to admit. He developed security software for eight years at CorreLog, where he authored the zDefender and SyslogDefender products which were acquired by BMC. He is currently the Chief Development Officer for Cloud Compiling and also does freelance projectoriented development.

He has a PhD in certificate technology from the University of Hard Knocks.

The University of Certificate Hard Knocks

Windows product

- My introduction to the nitty-gritty of certificates
- Implements both client and server ends of TLS protocol
- Built using OpenSSL
 - Open Source "lightly" documented forces one to learn as one goes
 - Result can be "the most dangerous code in the world" <u>https://bit.ly/2Djr76W</u>

z/OS product

- Built using IBM z/OS System SSL ("GSK")
- Designed to force you not to make mistakes
- Highly recommended

Why Certificates?

"Just a file"

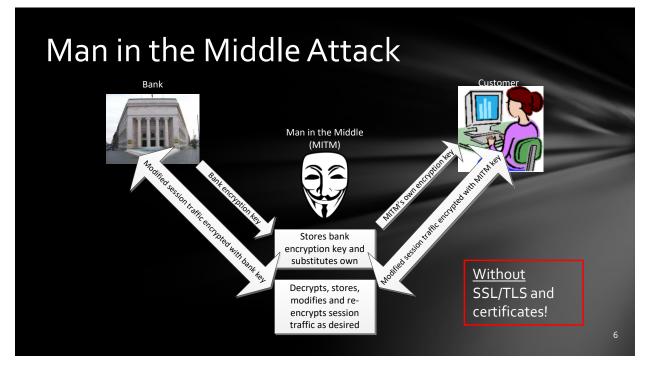
Automate security for remote connections

- Authentication: is this site really who it says it is?
- Encryption for data traffic
- For Web, FTP, TN3270 and potentially any similar connection

Authenticate users: is she really who she says she is?

Authenticate e-mail: is this e-mail really from the supposed sender, and how do I know it has not been altered?

Guarantee that software has not been tampered with since it left the publisher



Brief history of SSL and TLS ("SSL/TLS")

1994: Netscape, first company to commercialize the Internet, perceives browser communication not secure enough for e-commerce

1994: Develops Secure Sockets Layer (SSL) version 1 (never released)

1995 and 1996: SSL versions 2 and 3 (both now deprecated)

1998: Netscape crashes and burns in Microsoft browser wars

1999: SSL v3 becomes IETF Transport Layer Security (TLS) v1.0

TLS now at Versions 1.2 and 1.3. TLS 1.3 was defined in RFC 8446 in August 2018.

Certificates are a fundamental component of SSL/TLS

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100 MPH review of underlying technologies

• With links for additional reference

Details of the certificate protocol flow

100 MPH introduction to some advanced features

With links for additional reference

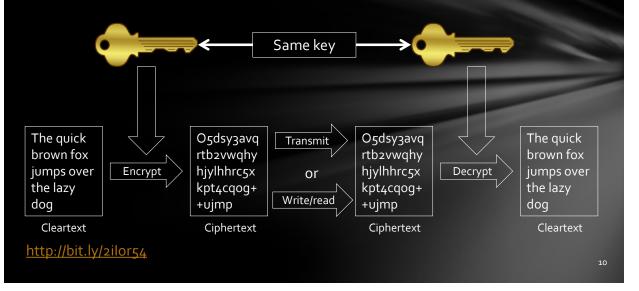


Netscape

AlphaCoders.com

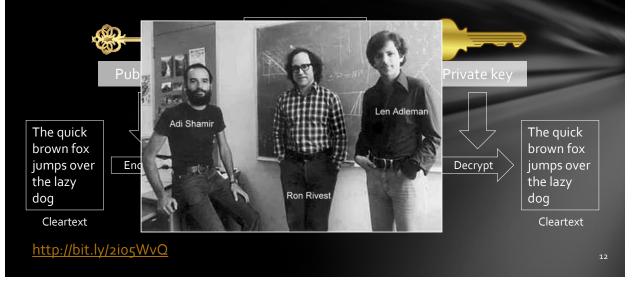


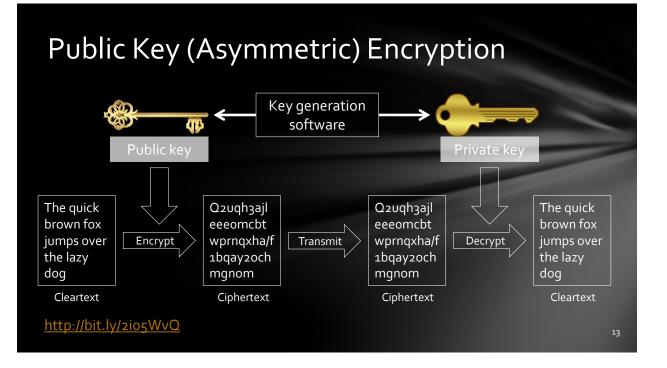
Secret (Symmetric) Key Encryption



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Public Key (Asymmetric) Encryption





My Public Key

Public-Key: (2048 bit) Modulus:

00:ad:3d:3a:cf:fd:39:8f:b0:d9:6d:8e:27:ad:37: c7:74:a2:b3:7a:05:b0:de:f9:06:96:f7:c6:a1:16: d5:2b:39:28:30:d2:63:3c:96:f5:3e:d1:9f:9b:9a: 1f:3e:29:71:be:7d:6b:c3:a3:90:de:ce:41:b0:e8: 5d:fe:ce:05:0d:d5:55:7f:fa:58:df:3b:5b:25:98: 8e:cb:c2:d1:6e:0d:be:44:88:87:9f:b1:a0:cf:de: ae:7d:e3:fd:d1:81:64:2b:48:f1:7a:83:d7:e9:66: 9f:32:3a:9a:26:d5:41:50:3e:8a:a4:9c:18:9a:c1: 21:ea:9b:b5:23:b1:57:27:55:e0:85:a0:d6:0e:c4: 3b:ea:8e:03:b7:4e:28:e0:c8:57:de:db:fe:a4:dc: 32:11:09:aa:d8:6d:04:e0:f6:d5:e2:08:c4:87:30: 29:3a:bd:0f:2b:45:7d:b8:6e:8c:71:22:ff:8c:3c: 68:7d:64:87:f7:87:a5:66:2c:d2:71:e2:97:84:48: 26:82:58:e4:0b:d6:59:e3:57:0a:07:24:77:e3:39: 3a:07:04:f6:ac:23:e1:33:28:ba:f3:5b:7c:df:91: 27:a4:79:a1:e5:6c:e9:c7:23:74:81:a7:cc:7f:75: c4:9e:d4:7e:27:af:23:9f:32:87:2e:f1:87:e7:38: 0f:31

Exponent: 65537 (0x10001)

But absolutely imperative to keep that private key private!

Public Key Encryption

What are the big problems?

Keys are HUGE! As you saw on the previous slide

- 4096 bits is over 1200 decimal digits
- Essence of public key is the difficulty of factoring huge numbers

Need a program to generate a pair of keys

Key management (are you noticing a theme?)

Unidirectional

Very slow

Don't say Mills said Public Key was no good – we will see how certificates solve all of these problems

Digital Hash

Function that takes a possibly very long "message" and returns a relatively short fixedlength binary value

Must be relatively easy to calculate

Deterministic: same message yields same hash

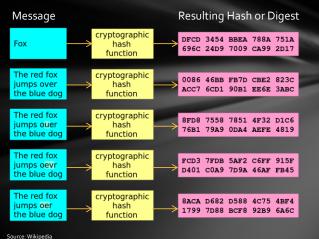
Possible but <u>highly</u> unlikely two slightly different messages have same hash

Almost impossible to construct a message with a predictable hash

Same hash = same message

Also called message digest

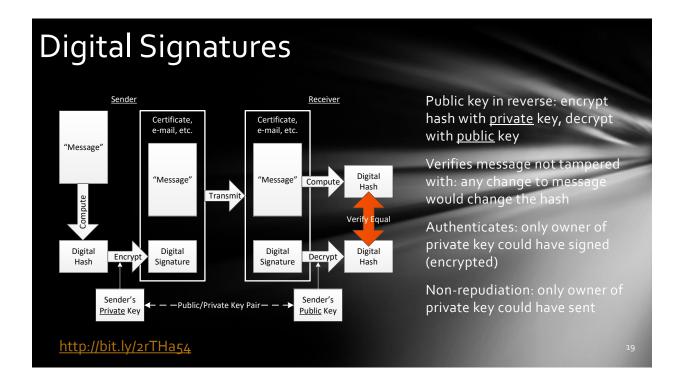
Examples: SHA-2 (MD5, SHA-1)



Source: Wikipedia By User:Jorge Stolfi based on Image:Hash_function.svg by Helix84 - Original work for Wikipedia, Public Domain, https://commons.wikimedia.org/wi/index.php?curid=s2s90240

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http://bit.ly/2mNglor



Digital Signatures and Trust

You receive a "message" (could be anything) purportedly from me with an attached digital hash encrypted with some private key. You also compute your own digital hash for the message.

If you can decrypt the attached digital hash with my public key (well known) and it is the same as the digital hash you computed, then the message is from me and is unaltered. If you trust me then you can trust the message.

By extension, if the message contains a public key, then you can trust any message signed with that public/private key pair.

This – the "chain of trust" – is the essence of certificates.

Certificate Authority

Company or group within a company that issues certificates

Uses self-created "root" certificate to sign them

May be well-known CA

- IdenTrust (owned by 8 large banks)
- DigiCert (20% market share)
- Comodo Sectigo (17% market share)
- Symantec (acquired Verisign*; CA business sold to DigiCert)
- GoDaddy
- GlobalSign
- Entrust*
- Let's Encrypt Free! Highly automated. Transparent (public log). Ninety days only!

Or department or individual

Well-known CA required for public-use SSL/TLS

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*CA Certificates that formerly shipped with RACF. IBM also shipped Thawte, acquired by Verisign (acquired by Symantec, acquired by DigiCert).

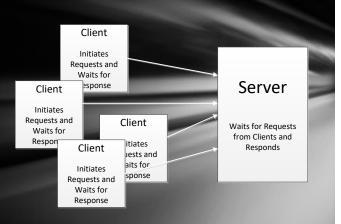
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Client and Server

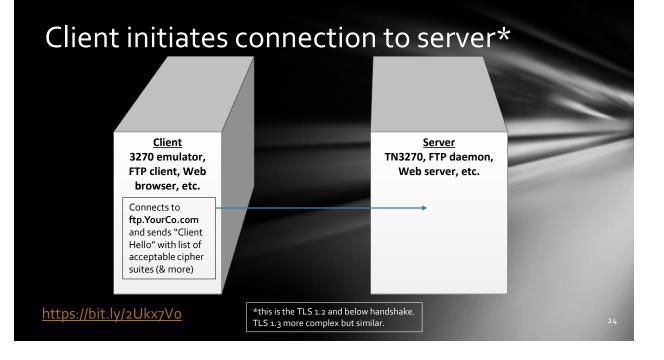
Nothing to do with color or size of boxes

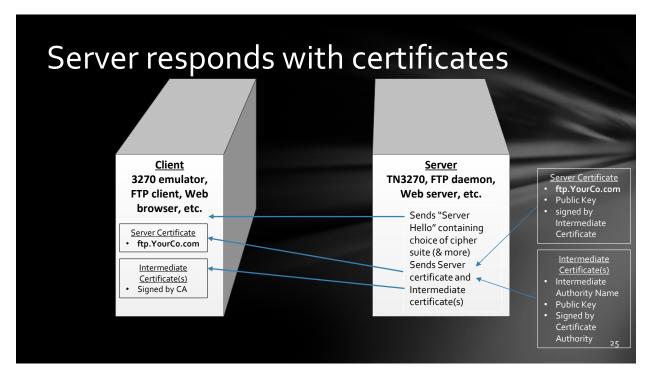
Often software, not hardware

The predominant architecture for complex applications (Web browser, FTP, e-mail, 3270 emulation)









What's in a Certificate?

The URL of the server for which it was issued	
The name of the certificate that signed this one	
MUST be unique within CA	
Start and end date and time	_
Half of this certificate's key pair	1000
Attests to the authenticity of this certificate	
	The name of the certificate that signed this one MUST be unique within CA Start and end date and time Half of this certificate's key pair

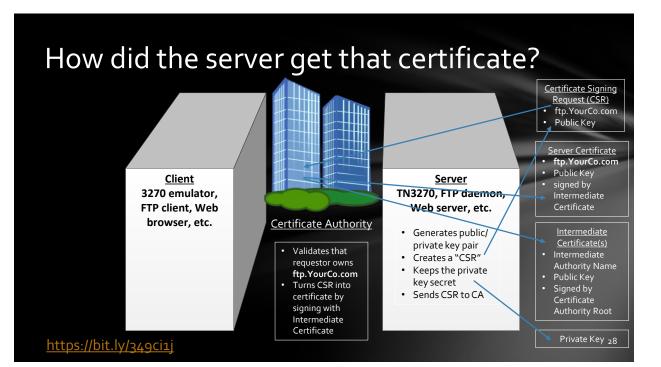
What's Never in a Certificate?

Private Key

Sometimes *packaged with* the certificate but never *part of* the certificate

https://bit.ly/3m8orXH

Formatted certificate con	tent
Label: CZAGENT_Nov2017_3 Trusted: Yes Version: 3 Serial number: 21 Issuer name: Charles Mills Consulting, LLC charlesm@mcn.org US	 Serial number "Common Name" (CN) of issuing CA Common Name of Subject
California Charles Mills Consulting, LLC Subject name: CZAGENT_NOv2017_3 charlesm@mcn.org US California Charles Mills Consulting, LLC Effective date: 2017/11/06	Validity dates Encryption algorithms Certificate Public Key
Expiration date: 2018/11/06 Signature algorithm: sha512WithRsaEncryption Public key algorithm: rsaEncryption Public key size: 2048 Public key: 30 82 01 0A 02 82 01 01 00 C1 56 C9 80 74 D7 EB 	Signature (not shown) Certificate Private Key may be <i>packaged with the</i> certificate but is
A2 42 5A A0 9F 7E 9E 3F 61 02 03 01 00 01 Above formatted display produced by IBM System SSL utility gskkyman. Get a similar display with the OpenSSL utility, which you can freely download and run on your desktop. <u>https://bit.ly/3qh95pO</u>	never <i>part of</i> the certificate. <i>Always</i> safe to transmit the certificate itself. 27



Why Intermediate Certificates?

The compromise of a CA root key would render root and all certificates issued by CA untrustworthy – a disaster!

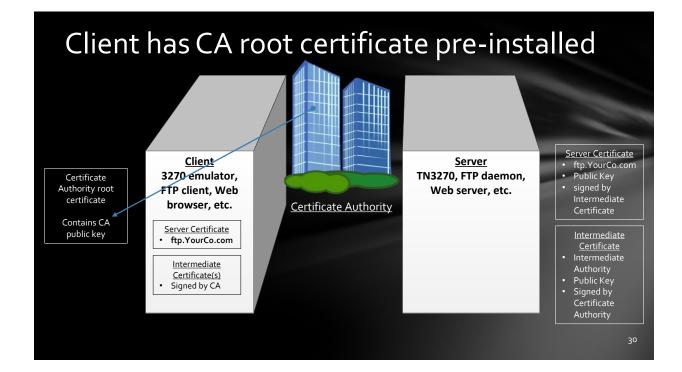
Certificate Authorities store their root keys off-line to help prevent compromise

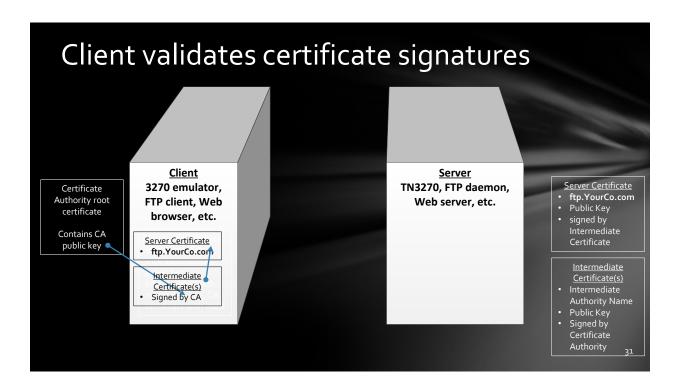
They use medium-term intermediate certificates – signed by their root certificate – to issue end-user certificates

Intermediate certificates are signed by the root certificate: "Chain of trust"

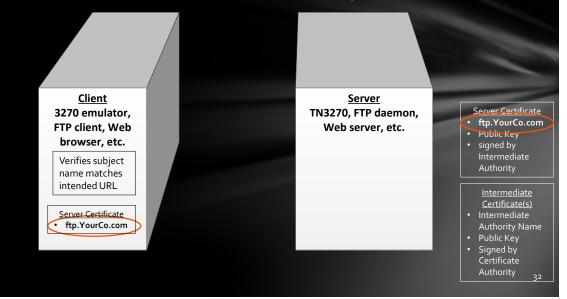
Owner's name		
Owner's public key		
Issuer's (CA's) name	reference	
Issuer's signature	Intermediate Cert	tificate
	Owner's (CA's) name	
sign	Owner's public key	
	Issuer's (root CA's) referen	nce
	Issuer's signature	*
	sign I	Root CA's name
		Root CA's public key
	self-sign	Root CA's signature
	Jen Jign	

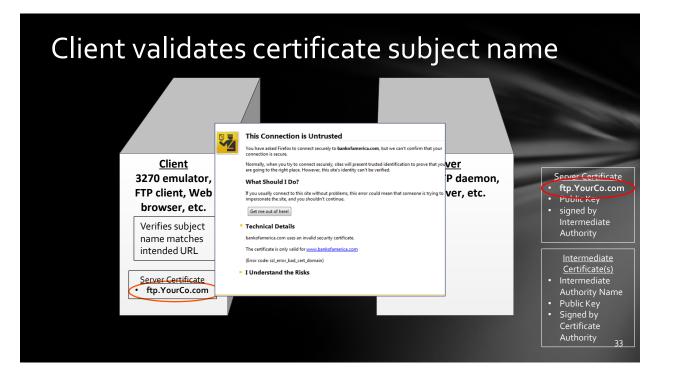
Source: Wikipedia By Yanpas - Own work, CC BY-SA 4.0, https://commons.wikimedia.org/w/index.php?curid=46369922





Client validates certificate subject name





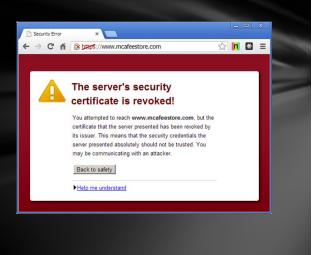
Certificate Revocation

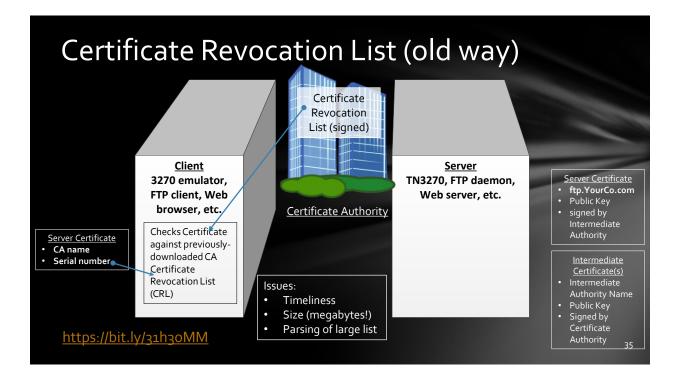
Why would a certificate be revoked?

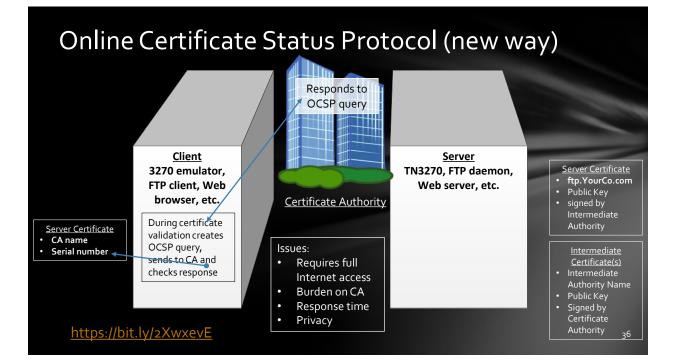
- Issued in error
- Key compromised
- CA root key compromised (disaster!)

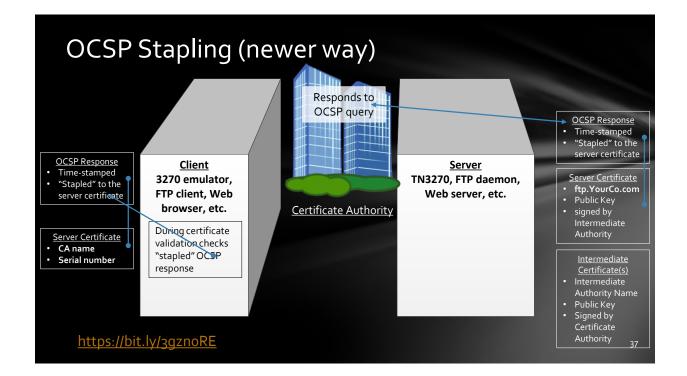
Clients should check for server certificate revocation

• Some clients do not, and some users ignore the error – bad idea!

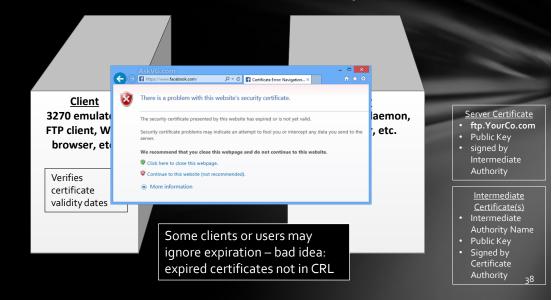


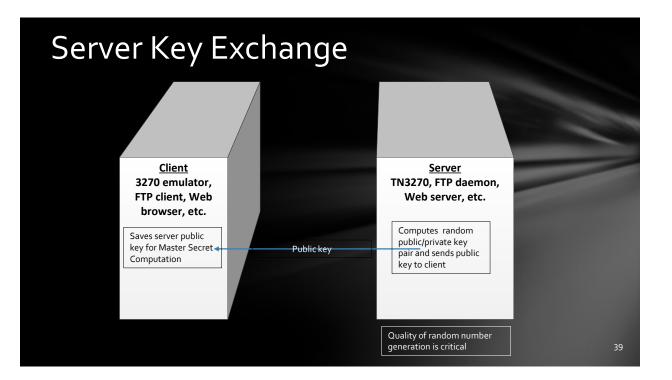


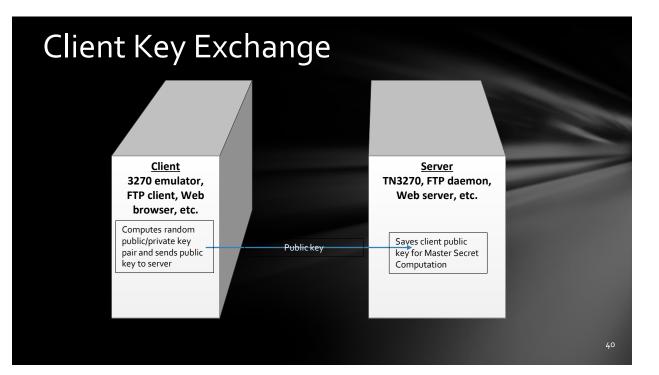




Client checks certificate validity dates







Client and Server Compute Master Secret

Client and Server each compute a "Premaster Secret"

- Computed from random numbers in Client Hello, Server Hello, its own private key, and the partner's public key
- Both parties perform the same computation and should get the same result (even though different inputs!)
- Length varies depending on cipher suite

Client and Server each derive the same "Master Secret" from the Premaster Secret

• Always 48 bits

Up to Six session keys derived from master secret Client Write Encryption Key Client Write Initialization Vector Key (used only for certain ciphers)

Client Write MAC Key (used for message authentication)

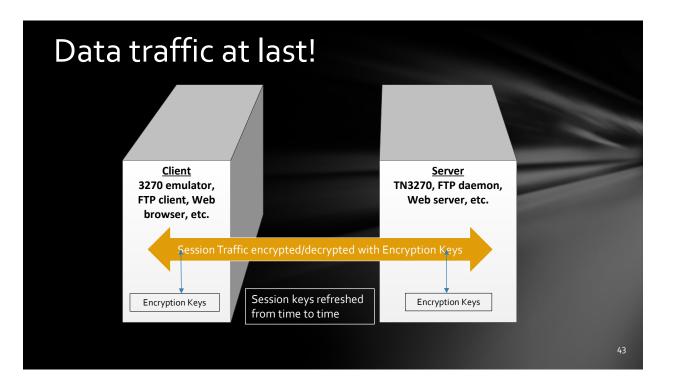
Server Write Encryption Key

Server Write Initialization Vector Key (used only for certain ciphers)

Server Write MAC Key (used for message authentication)

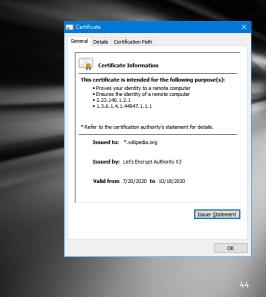
https://bit.ly/3ho66OB

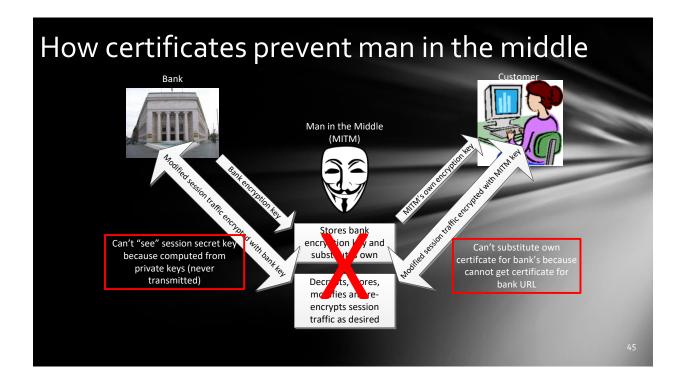
Master Secret



Certificates Solve the Crucial Problems

- Authentication
- Encryption
- Secure key delivery
- Automation of key delivery
- **Bi-directionality**
- Speed
- Man in the Middle attack





Certificate Issues

Complexity

Certificate management

Especially expiration

Key management

- Keeping private key private
- But not losing them!

CA Root Certificates and Trust

Certificate Authority Issues

- Sloppiness, fraud?
- Repressive government pressures CA to facilitate Man-in-the-Middle
 - Dutch CA DigiNotar hacked; fraudulent Google.com certificate used for Man-in-the-Middle interception of Iranian citizens https://bit.ly/ahNaban
- Name validation by CA
 - Requirement for CA to validate URLs at odds with modern certificate volumes
 - In March of 2017, Google announced Chrome would stop honoring Symantec certificates for (among other things) sloppiness in validating certificate names https://bit/bi20G51xd
 - Death penalty! Symantec sold CA business to DigiCert
- Any CA can issue a certificate for any site!





Self-signed certificates

Misunderstood concept

Self-signing is not inherently bad – all CA root certificates are self-signed

Means the certificate signs <u>itself</u>, not that the company that issued the certificate is its own CA

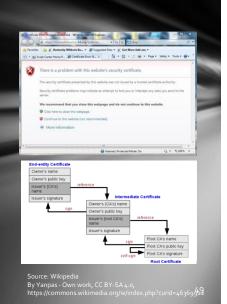
Generally frowned upon for end-point certificates

Provide encryption

Provide authentication only if pre-installed on client

Nothing wrong with your company being its own CA

- Saves money, time and trouble
- Works only for internal clients external users do not have CA root certificate
- Possibly more secure to control it all yourself



Alternative Names

Certificates support multiple "subject alternative names" (SANs) in addition to the main "common name"

Thus one certificate could be valid for YourCo.com, MyCo.com and HerCo.com

Using an Alternative Name for the server URL is now preferred to Common Name (RFC 2818)

Sometimes called a Multi-Domain or SAN Certificate

CA's charge more for multiple names but that is a business issue, not a technical issue

http://bit.ly/2B8AL4Z

Subject Name Wildcards

Certificates support wildcard subject names (Common or Alternative)

Asterisk may be last or only character of leftmost subdomain name: *.YourCo.Com or w*.YourCo.Com

• Or last dotted address octet: 192.168.17.* (infrequent)

One certificate for www.YourCo.com, ftp.YourCo.com and mail.YourCo.com

Certificate	×
neral Details Certification Path	
Certificate Information	
This certificate is intended for the following	g purpose(s):
 Ensures the identity of a remote computer 2.23.140.1.2.2 	
-	
Issued to: *.google.com	
Issued by: GTS CA 101	
Valid from 7/7/2020 to 9/29/2020	
	Issuer <u>S</u> tatement
	OK

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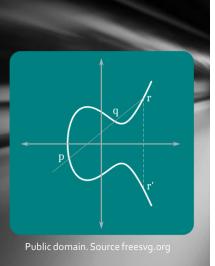
Elliptic Curve Encryption

The problem with RSA encryption

- Principle is that multiplication is fast; factoring is slow
- As computers have gotten faster we have compensated by going to larger and larger RSA keys
- Problem is that the larger the key, the less the difference in time between multiplication and factoring – so diminishing returns

Elliptic Curve Encryption too complex for one slide

- Relatively fast to compute a transformation based on an elliptic curve
- Very slow to reverse that transformation
- Smaller keys give security equivalent to large RSA keys
- Time ratio constant for larger keys



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Diffie-Hellman and Forward Secrecy

Perfect Forward Secrecy

- RSA key exchange uses certificate private key to derive encryption key
- Suppose intruder stole server certificate private key
- You could just re-issue the certificate with a new key
- But suppose the intruder had recorded earlier session traffic
 He could now decode it all with his stolen key
- TLS 1.3 prevents by requiring "perfect forward secrecy"

Ephemeral Diffie-Hellman (DHE) Key Exchange

- Client and Server separately compute premaster secret from partner's public key and own private key (+ exchanged random numbers)
- They arrive at identical result, but intruder has neither private key and cannot
- Key is "ephemeral" and used for only the one session
- Hence intruder cannot decode using stolen certificate key

https://bit.ly/3zZuDMm



Image credit: Chuck Painter/Stanford News Serv

Client Certificates

Server certificate authenticates server identity and provides for encryption

Client certificate authenticates client identity only

- Does not provide for or configure encryption
- Must be CA-signed or else pre-installed on server

An alternative to passwords

Good choice if relatively small number of clients, over which you have control

Good for branch offices, not for customers

Server makes protocol request for certificate from client, so configuration is a server option

FTP Example (server-side): SECURE_LOGIN VERIFY_USER

Validation protocol similar to server certificate

Code signing with certificates

Verifies that software is authentic

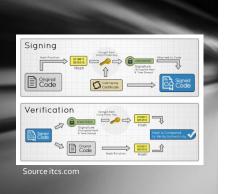
Does not prove that code is good, merely authentic!

Verifies software has not been altered/tampered with

Requires special code-signing software

May be CA-signed or software-vendor signed

- **Time-stamping**
- Allows for fact that certificate may expire after software is published but before it is installed



Constraints and Key Usage

Basic constraint

CA key or not

Key usage

- Signatures
- Etc.

Extended key usage

- Server
- Client
- Code signing
- Email
- Etc.

http://bit.ly/2B7lZvy

X509v3 extensions:

X509v3 Basic Constraints:

CA:FALSE

X509v3 Extended Key Usage:

TLS Web Server Authentication, TLS Web Client Authentication

Summary

Why certificates?

100 MPH review of underlying technologies

With links for additional reference

Details of the certificate protocol flow

100 MPH introduction to some advanced features

With links for additional reference

More questions? charlesm@mcn.org



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